

Taming Metadata Burstiness of **HPC Jobs Through Application**level QoS Control

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PADLL IS AN

SYSTEM AGNOSTIC

STORAGE MIDDELWARE

THAT ENABLES QOS

CONTROL IN HPC

STORAGE SYSTEMS

INESC TEC and University of Minho AIST TACC and UT Austin Intel

PROBLEM STATEMENT

EFFECTIVELY ENSURING STORAGE QOS GUARANTEES IN LARGE-SCALE HPC SYSTEMS IS NOT TRIVIAL:

Manual intervention

- In HPC facilities, sysadmins manually stop jobs with aggressive I/O behavior
- Reactive approach that is only triggered when concurrent jobs were already harmed

Intrusive to I/O layers

- Existing solutions are tightly coupled to core layers of the HPC stack (e.g., Parallel File System (PFS), job scheduler)
- Require profound system refactoring

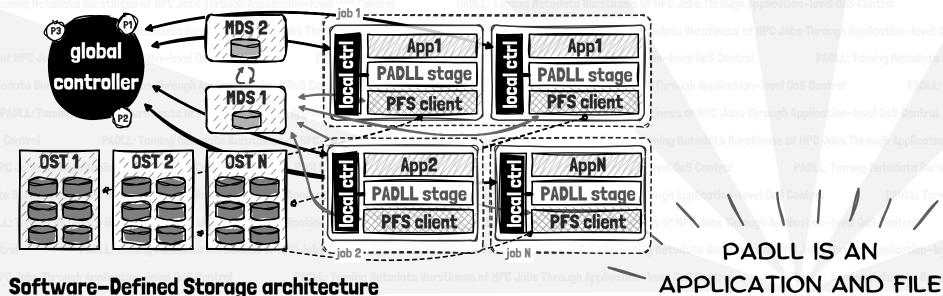
Partial visibility and I/O control

- Compute node-level solutions actuate in isolation (i.e., agnostic of other jobs)
- multiple jobs that compete for shared storage

Metadata remains overlooked

- Existing proposals mainly focus on achieving QoS over data workflows
- Metadata operations of a single job can saturate the PFS metadata resources

22 PADLL STORAGE MIDDLEWARE



Software-Defined Storage architecture

• PADLL is organized in multiple data plane stages that differentiate and rate limit I/O workflows, and a hierarchical control plane that manages all stages to ensure storage QoS policies

Application and PFS agnostic

- PADLL sits between applications and the PFS, and does not require changing any core layer of the HPC I/O stack (LD_PRELOAD)
- Compatible with POSIX-compliant storage systems

Fine-grained I/O control

• Classifies, differentiates, and enforces I/O requests with different levels of granularity (e.g., operation type, class, and job)

Global visibility

 Coordinated control of all I/O workflows destined towards the PFS preventing I/O contention and unfair usage of shared resources

• Unable to coordinate the I/O generated from

Uniform rate distribution

• Jobs are throttled with a fixed rate throughout their execution, regardless of their size, duration, and workload

3 CONTROL ALGORITHMS

• Even distribution of resource shares

Priority-based rate distribution

- PFS resources are distributed based on a given priority
- Can lead to both under-provisioning (e.g., leftover I/O resources) and over-provisioning (e.g., resource shares larger than needed)

Proportional sharing (psharing)

- Traditional max-min fair share control algorithm (feedback loop) that enforces per-job rate reservations (IOFlow, Retro, PAIO)
- Suited for workloads with sustained I/O load, but suboptimal under volatile workloads (over-provisioning)

Proportional sharing without false resource allocation (psfa)

- New max-min fair share algorithm (feedback loop) that ensures storage QoS under volatile workloads
- Assigns resource shares based on the <u>actual I/O usage</u> of each job and their respective metadata demands

引 EXPERIMENTAL TESTBED

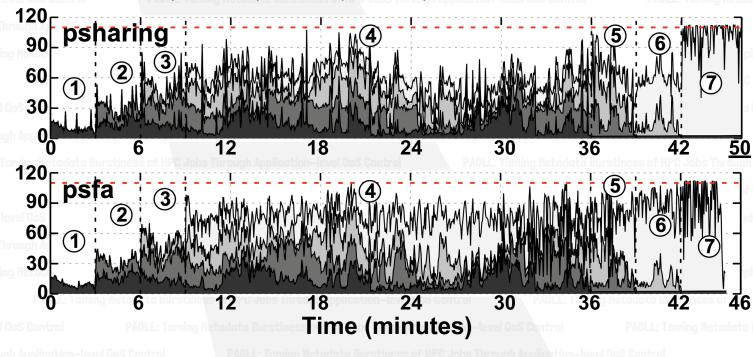
• Limit overall metadata load in the PFS (110 kops/s), while assigning different I/O priorities to jobs

Experimental environment

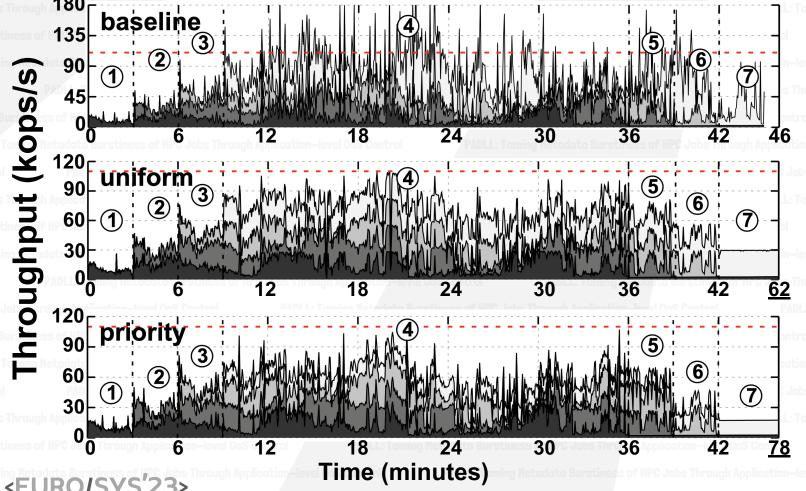
- Per-job QoS control in the <u>ABCI supercomputer</u>, hosted by AIST
- Metadata traces from a production Lustre file system
- Experiments include 7 phases, each marking when a job enters or leaves the system

<u>Setups</u>

- Baseline 4 jobs with different loads (15%, 20%, 20%, 45%)
- Uniform all jobs set to 27.5 kops/s
- Priority, psharing, and psfa all jobs are assigned with different rates (40) kops/s, 25 kops/s, 30 kops/s, 15 kops/s)



5 PER-JOB METADATA CONTROL



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